



ECML
PKDD
2020



中国科学院计算技术研究所
INSTITUTE OF COMPUTING TECHNOLOGY, CHINESE ACADEMY OF SCIENCES



SPECGREEDY

Unified Dense Subgraph Detection

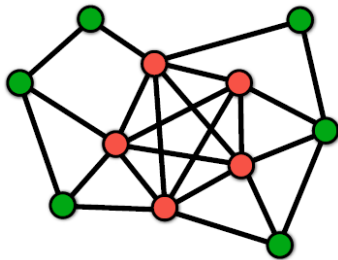
Wenjie Feng⁺, Shenhua Liu⁺, Danai Koutra[#],
Huawei Shen⁺, Xueqi Cheng⁺

⁺*Institute of Computing Technology, ICT, CAS*

[#]*University of Michigan, Ann Arbor*

Motivation: Dense Subgraph Detection

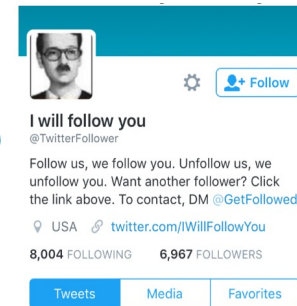
- Ubiquitous applications



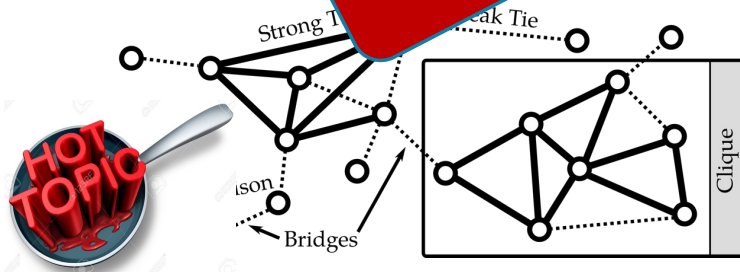
Densest subgraph
[Charikar M'06]



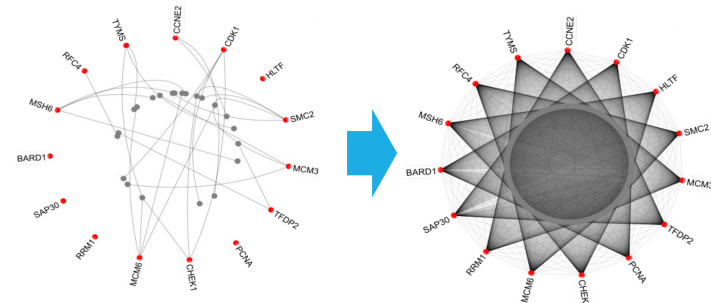
Community detection
[Katz A. et al'18]



Fraud detection
[B. Hooi et al'18]



News hot topics discovery
[LY, Chu et al'12]

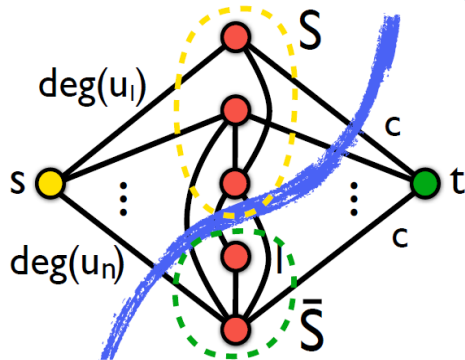


Contrast dense subgraph
[Wong S. W e al'18]

Motivation: Densest Subgraph Detection

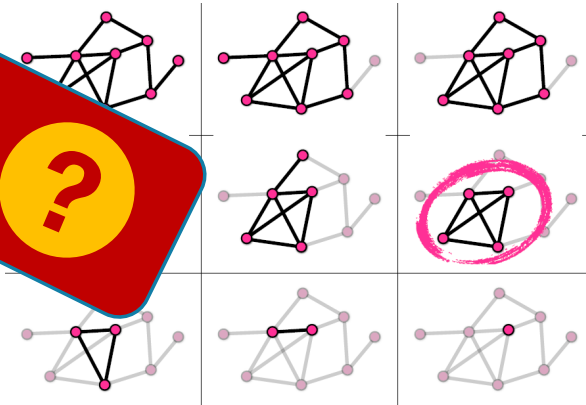
- **Polynomial-time methods**

- **Exact solution:** maximum flows based algo.
 - $O(|U|^2 + |U| \cdot |E|)$
- **Approx. solution:** Greedy Peeling algo. ($1/2$ -optimal)
 - $O(|E| \cdot \log|U|)$



Max-flow exact DS
[Goldberg'84]

faster algorithm



Greedy approx. DS
[Charikar,M.'00]

OVERVIEW-METHOD

- **GENDS**: a unified problem framework

$$S^* = \arg \max_{S \subseteq V, |S| \geq 1} g(S; \mathbf{P}, \mathbf{Q}) = \arg \max_{\mathbf{x} \in \{0,1\}^n, |\mathbf{x}| \geq 1} \frac{\mathbf{x}^T \mathbf{P} \mathbf{x}}{\mathbf{x}^T \mathbf{Q} \mathbf{x}}$$

$$\mathbf{P} = \mathbf{A} + 2\mathbf{D}_c \quad \mathbf{Q} = \mathbf{A}' + \gamma \mathbf{I}$$

MinQuotientCut

[Chung RK. Fan'96]

SparseCutDS

[Miyauchi A. et al'18]

TempDS

[Wong S. W e al'18]

Fraudar

[B. Hooi et al'18]

Charikar

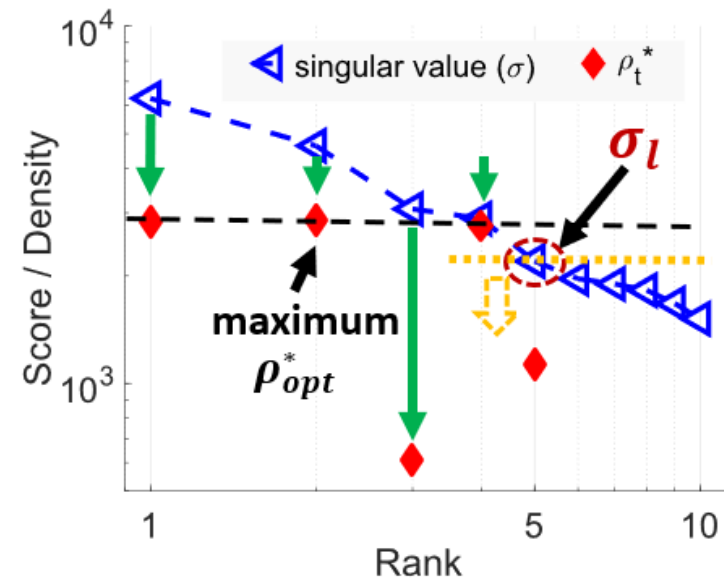
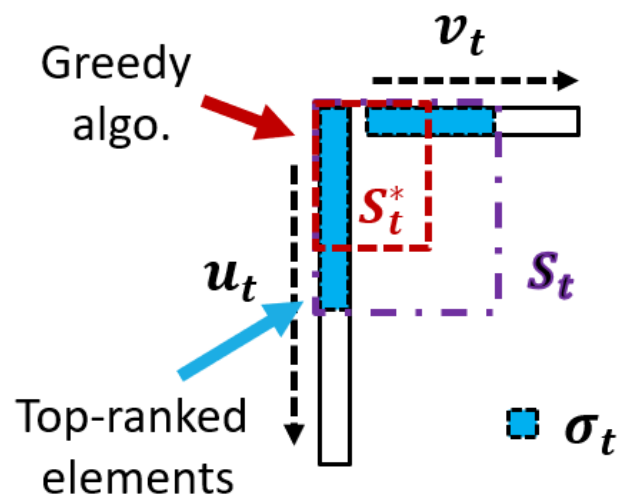
[Charikar M.'00]

Risk-averse DS

[C. E. Tsourakakis et al'19]

OVERVIEW-METHOD

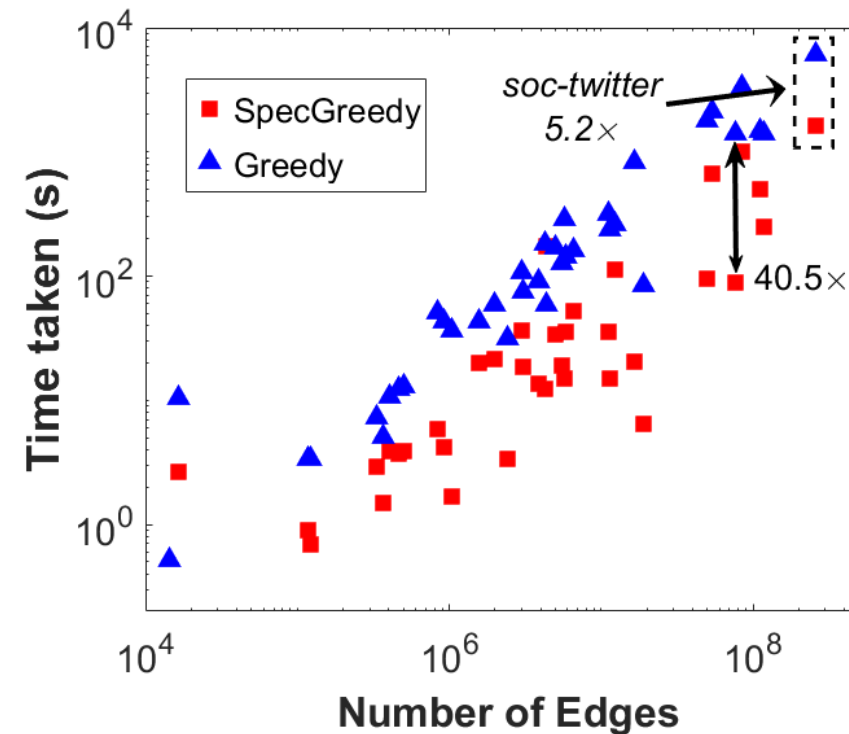
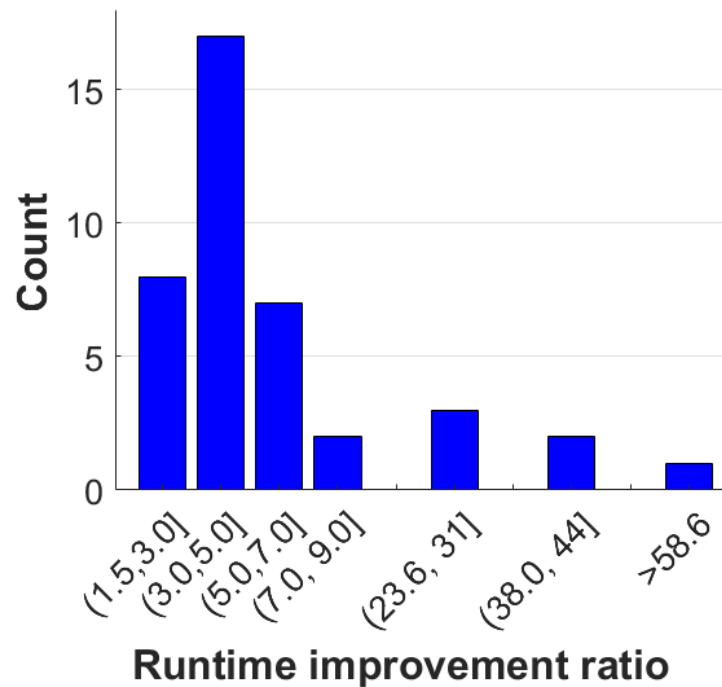
- **SPECGREEDY**: a unified detection algorithm



Graph spectral + Greedy approximation

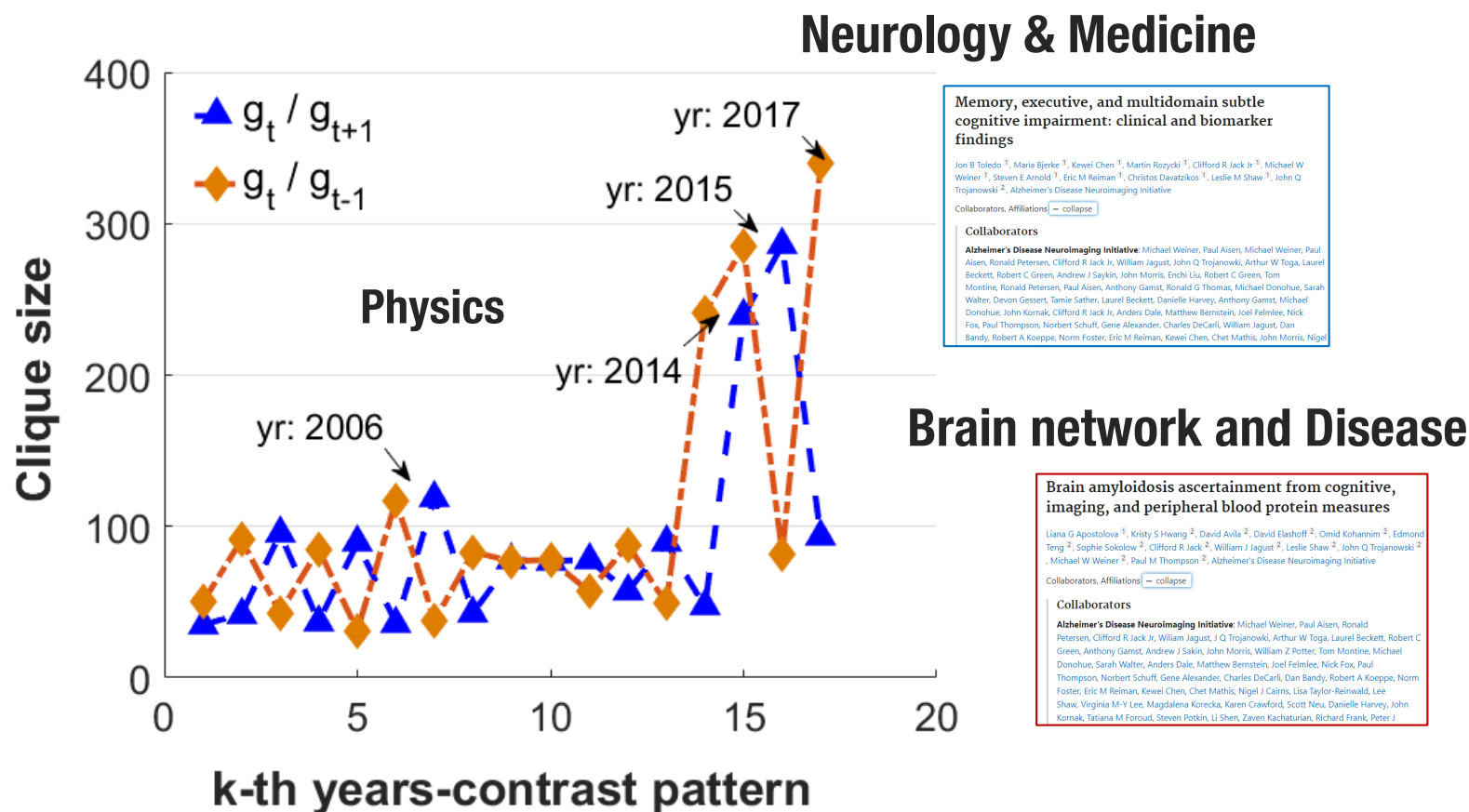
Linearly scalable with the graph size

OVERVIEW-PERFORMANCE



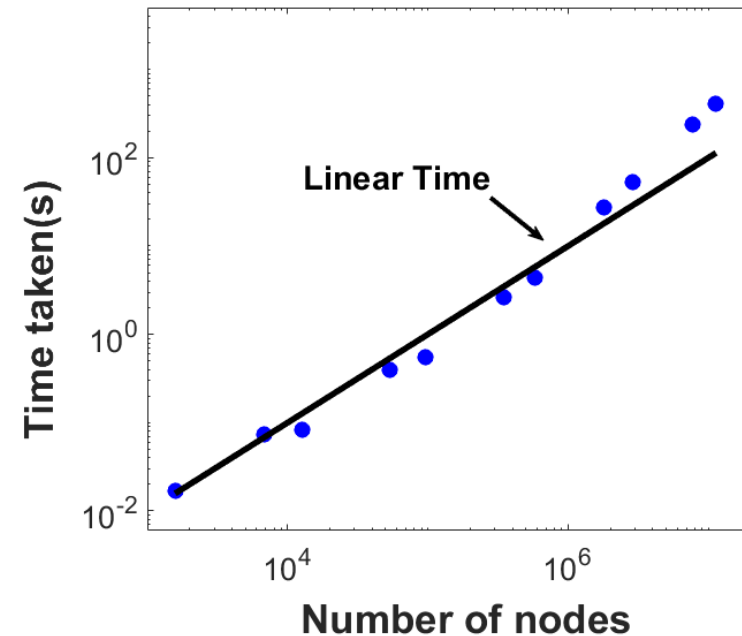
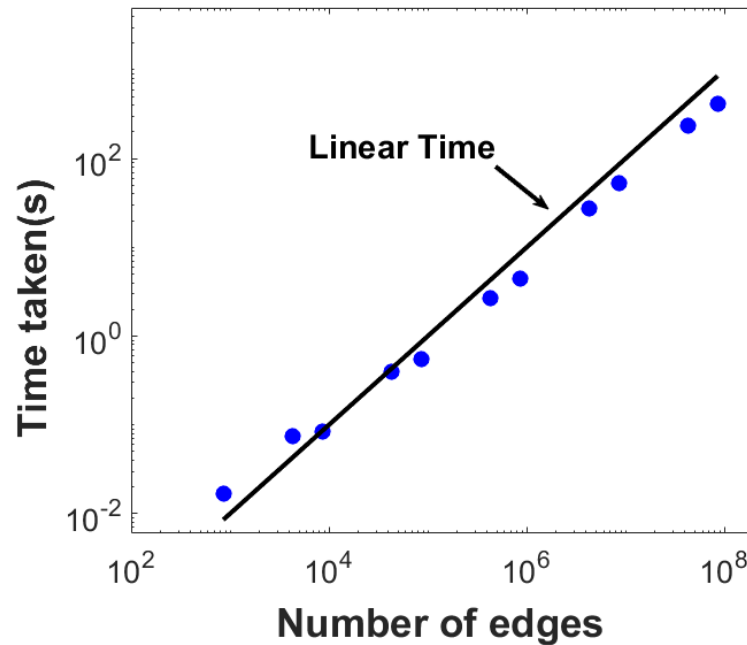
SPECGREEDY is faster than Greedy algo.

OVERVIEW-PERFORMANCE



SPECGREEDY is effective, e.g. contrast dense subgraphs

OVERVIEW-PERFORMANCE



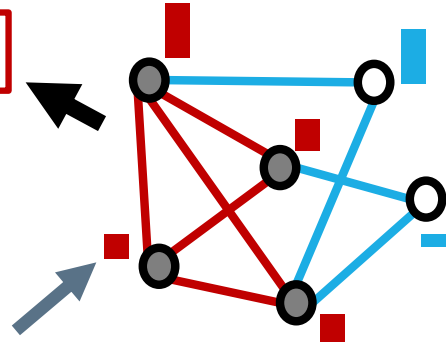
SPECGREEDY is scalable

Road Map

- Introduction
- Proposed Method:
 - Unified Formulation: **GENDS** <<
 - Algorithm: SPECGREEDY
- Experiments
- Conclusion

GENDS: Unified Formulation

- **GENDS: Generalized Dense Subgraph detection**
- **Given:**
 - $\mathcal{G} = (U, E)$ and its contrast $\mathcal{G}' = (U, E')$
 - matrices $\mathbf{P} = \mathbf{A} + 2\mathbf{D}_c$, $\mathbf{Q} = \mathbf{A}' + \gamma\mathbf{I}$ related to \mathcal{G} and \mathcal{G}' resp.
- **Find**
 - the optimal subset $\mathbf{S}^* \subseteq U$,

$$S^* = \arg \max_{S \subseteq V, |S| \geq 1} g(S; \mathbf{P}, \mathbf{Q}) = \arg \max_{\mathbf{x} \in \{0,1\}^n, |\mathbf{x}| \geq 1} \frac{\boxed{\mathbf{x}^T \mathbf{P} \mathbf{x}}}{\mathbf{x}^T \mathbf{Q} \mathbf{x}}$$


c_i : node weight

\mathbf{x} is the indicator vector for the node subset

Problems Correspondence

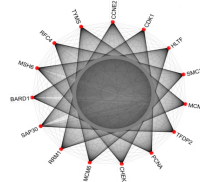
- GENDS and some related instantiation problems

$$c_i = -\frac{d_i}{2}$$

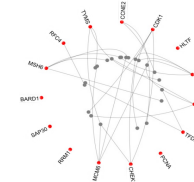
$$\mathbf{A}' = \mathbf{0}, \gamma = 1$$

	Method	matrix \mathbf{P}	matrix \mathbf{Q}
1	MinQuotientCut	$\mathbf{A} - \mathbf{D} = -\mathbf{L}$	\mathbf{I}
2	Charikar	\mathbf{A}	\mathbf{I}
3	Fraudar	$\mathbf{A} + 2 \mathbf{D}_w$	\mathbf{I}
4	SPARSECUTDS ¹	$\mathbf{A} - \frac{2 \cdot \alpha}{2\alpha+1} \mathbf{D}$	\mathbf{I}
5	TEMPDS	\mathbf{A}_t	$\mathbf{A}_{t-1} + 2 \mathbf{I} = \tilde{\mathbf{A}}_{t-1}$
6	Risk-averse DS	$\mathbf{A} + \lambda_1 \mathbf{I} = \tilde{\mathbf{A}}^+$	$\mathbf{A}^- + \lambda_2 \mathbf{I} = \tilde{\mathbf{A}}^-$
	GENDS ²	$\mathbf{A} + 2 \mathbf{D}_c$	$\mathbf{A}' + \gamma \mathbf{I} = \tilde{\mathbf{A}}'$

Graph at time t



Graph at time $t-1, \gamma = 2$



Theoretical Analysis

• 1. Graph spectral and Optimality

◦ Generalized *Rayleigh ratio* (quadratic optimization)

■ Positive residual graph \mathcal{G}_r : $\mathbf{A}_r = (\mathbf{P} - \mathbf{Q})^+$

■ $R(\mathbf{A}_r, \mathbf{x}) = \frac{\mathbf{x}^T \mathbf{A}_r \mathbf{x}}{\mathbf{x}^T \mathbf{x}}, \quad \mathbf{x} \in \mathbb{R}^n, \mathbf{x} \neq \mathbf{0}$

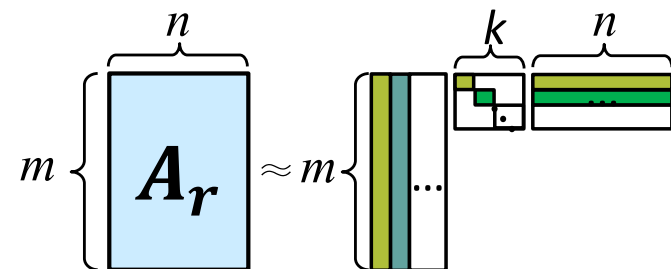
◦ *Rayleigh-Ritz* Theorem (real space)

■ $\mathbf{A}_r = \mathbf{U} \mathbf{\Lambda} \mathbf{U}^T, \mathbf{\Lambda} = \text{diag}(\lambda_1, \dots, \lambda_n)$

■ $\lambda_k = \max_{\mathbf{x} \neq \mathbf{0}, \mathbf{x} \in \mathcal{S}_{k-1}^\perp} R(\mathbf{A}_r, \mathbf{x}) = \max_{\|\mathbf{x}\|=1, \mathbf{x} \in \mathcal{S}_{k-1}^\perp} \mathbf{x}^T \mathbf{A}_r \mathbf{x} \implies \mathbf{x} = \mathbf{u}_k$

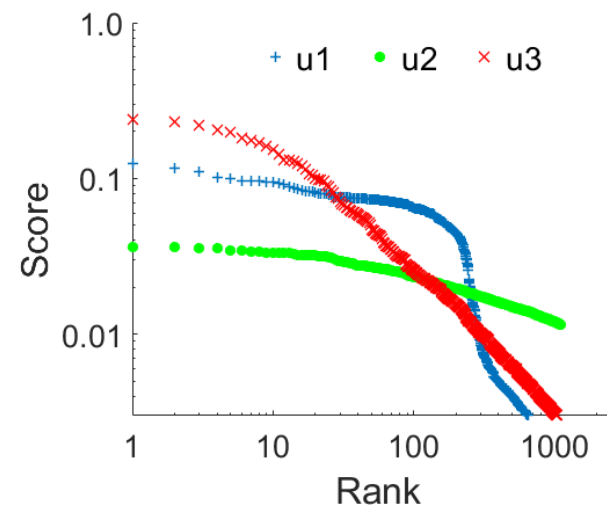
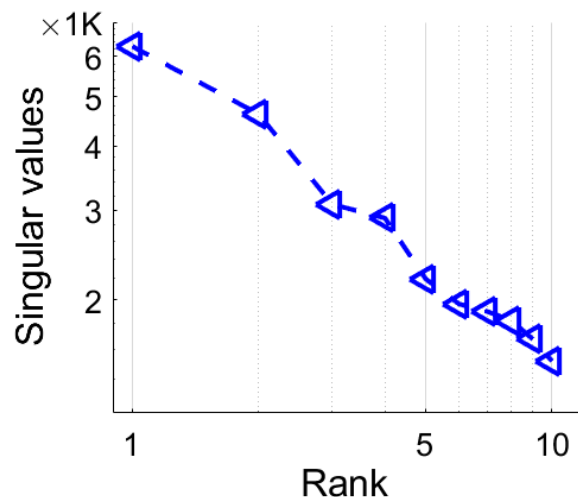
◦ $\mathbf{A}_r = \mathbf{U} \mathbf{\Sigma} \mathbf{V}^T, \mathbf{\Sigma} = \text{diag}(\sigma_1, \dots, \sigma_k)$

$$\sigma_1 \geq \dots \geq \sigma_k \geq 0$$



Theoretical Analysis (cont.)

- 2. Properties of real-world large graphs
 - Sparse connection
 - Skew distributions (e.g., power-law)
 - degree, eigen/singular value, network values, |bipartite core| etc.



Singular values and score of singular vector elem. distribution of **soc-twitter** network

Theoretical Analysis: Optimal Solution

- GENDS solution for \mathcal{G}_r

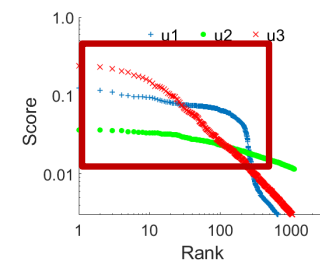
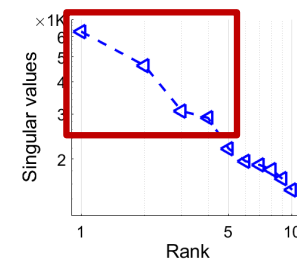
- $$S^* = \arg \max_{\mathbf{x} \in \{0,1\}^n, |\mathbf{x}| \geq 1} \frac{\mathbf{x}^T \mathbf{A}_r \mathbf{x}}{\mathbf{x}^T \mathbf{x}} = \arg \max_{|S| \geq 1} \frac{1}{|S|} \sum_{i=1}^n \sigma_i \left(\sum_{j \in S} u_{ij} \right) \left(\sum_{j \in S} v_{ij} \right)$$

- The optimal density of subgraph of S^* :

$$\rho_{opt} \leq \sigma_1$$

- Important parts:

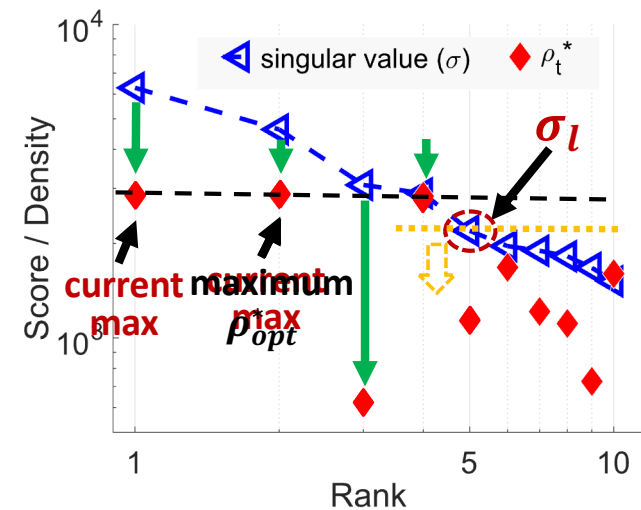
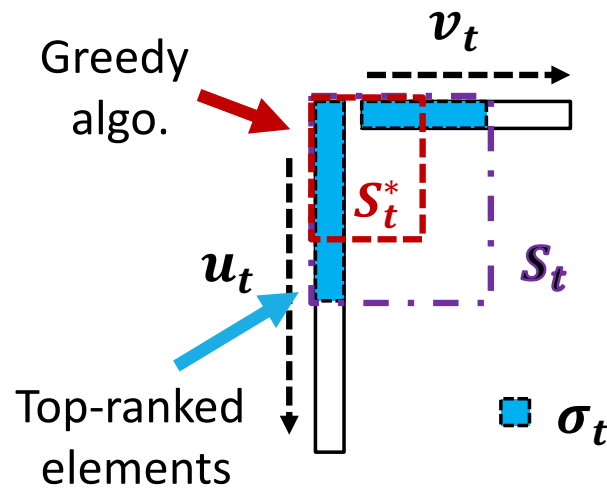
- top singular values
- top-ranked elements in singular vectors



Road Map

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 - **Algorithm: SPECGREEDY <<**
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Approximate Optimal Solution



- Approximation solution S_t^* :
 - g_t : subgraph for S_t of the trunc. t -th top singular vectors (u_t, v_t)
 - g_t^* : the densest subgraph of g_t with Greedy algo.
 - Density of g_t^* : $\rho_t^* \leq \sigma_t$
- S_{opt}^* approx. optimal solution for \mathcal{G}_r : $\sigma_1 > \dots > \rho_{opt}^* \geq \sigma_l \geq \dots$

SPECGREEDY: Unified Detection Algorithm

- Graph spectral + Greedy peeling approx.
- Main steps: (k : parameter for SVD approx.)
 - S0: construct positive residual graph \mathcal{G}_r
 - S1: A_r SVD (top- k low-rank approx.)
 - Repeat:
 - S2 {
 - S2-1: Construct candidate S_t for truncated (u_t, v_t)
 - S2-2: S_t^* : densest subgraph with greedy detection for S_t
 - S2-3: **Update** current optimal density ρ_{opt}^* with S_t^* if needed
 - Until $\rho_{opt}^* > \sigma_{t+1}$ or $t < k$
 - Return $\mathcal{G}_r(S_{opt}^*)$

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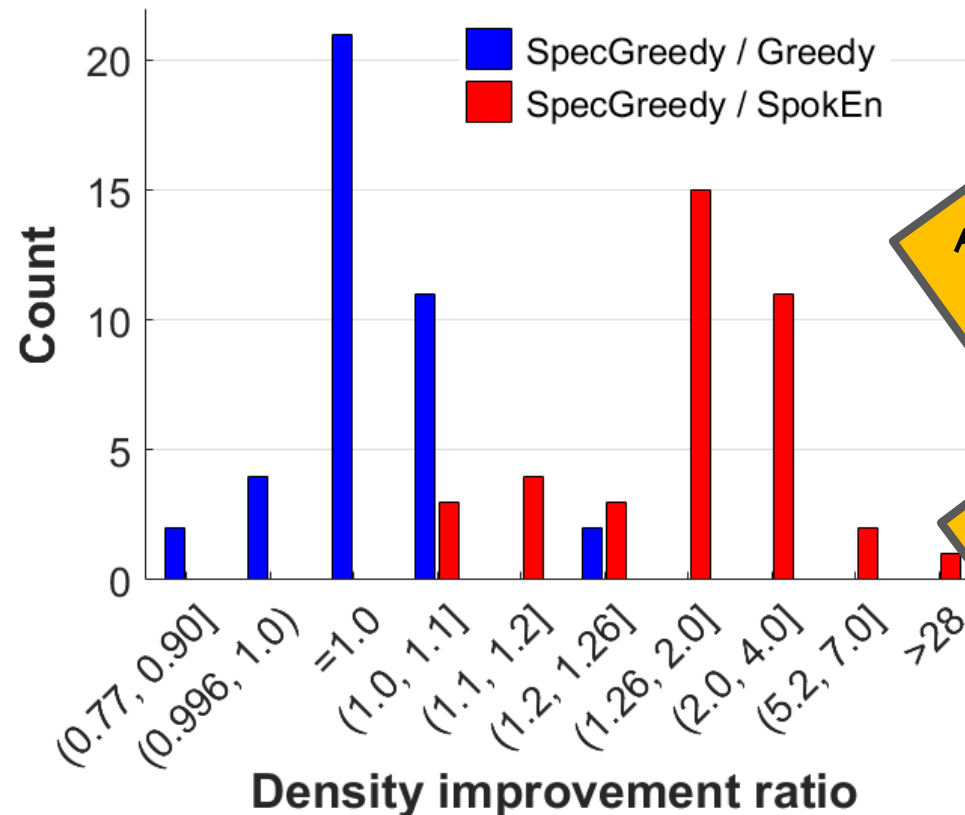
Experiments on 40 networks

- Baselines:
 - Greedy (Charikar)
 - SpokEn (SVD)
- Parameter: $k = 10$ for SPECGREEDY
- Datasets:
 - **40** real-world networks from **5** public dataset repos.
 - Stanford SNAP, Network Repository, AMiner scholar, ASU' social computing, KONECT Network Collection



Exp1. Quality-Speed Test

- The **density** of the detected densest subgraph

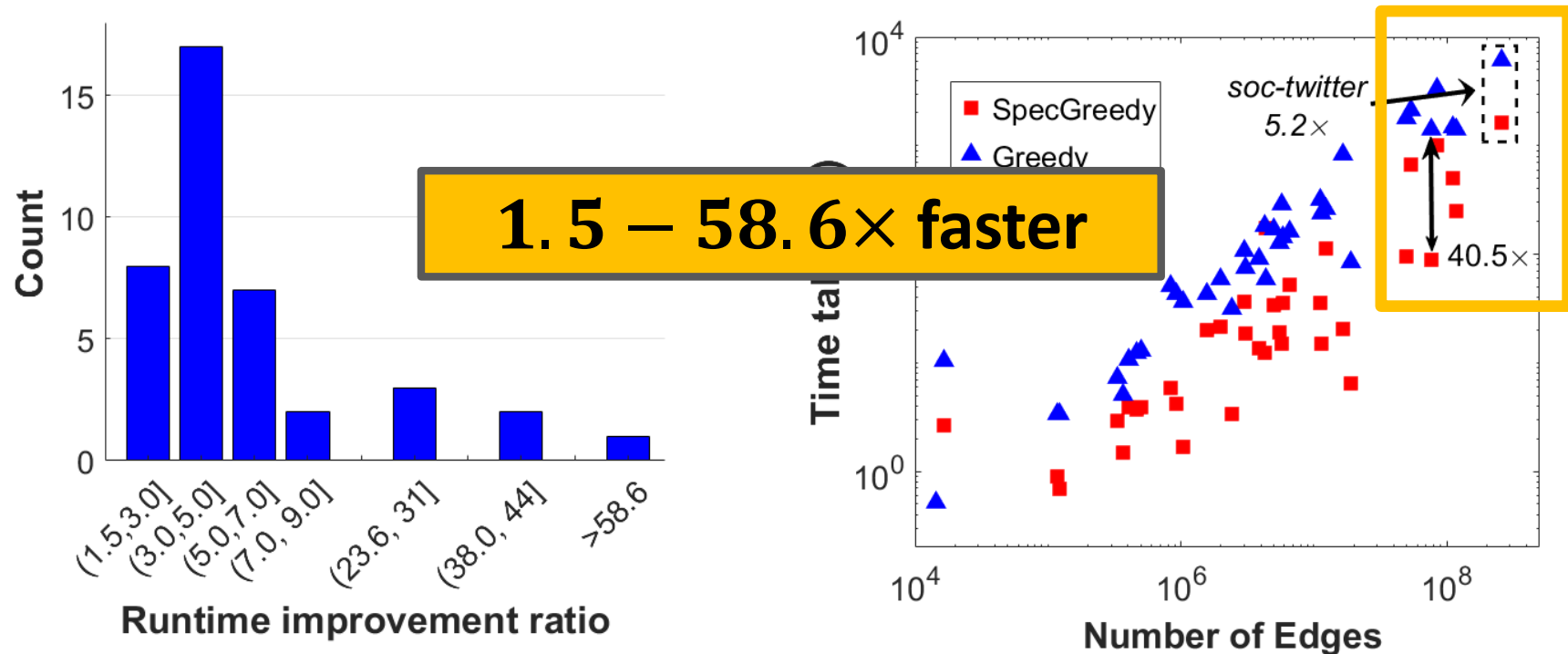


1.26× denser
than Greedy

77× denser
than SpokEn

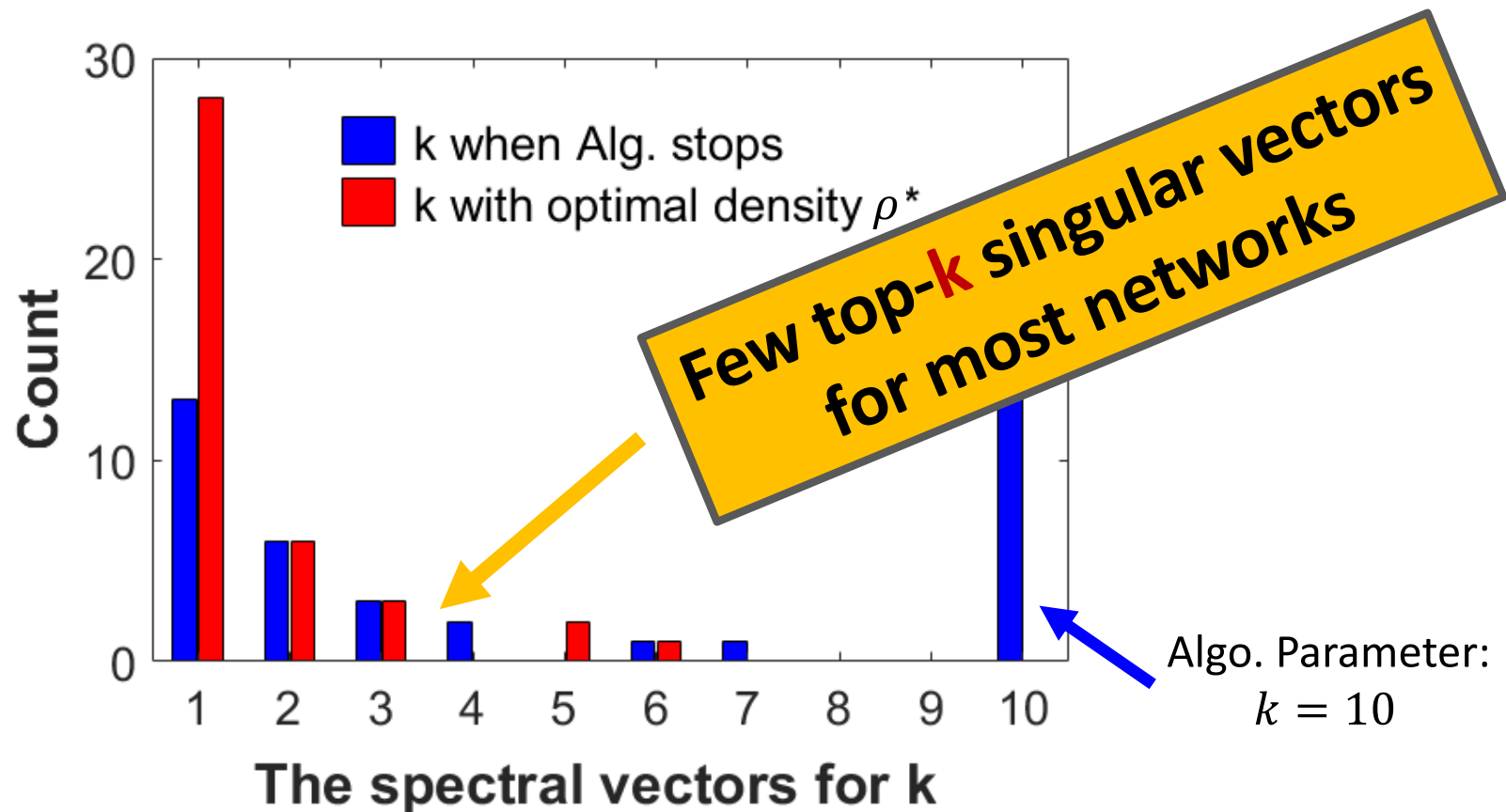
Exp1. Quality-Speed Test (cont.)

- The **speed-up** for detecting the densest subgraph
 - SpecGreedy vs. Greedy



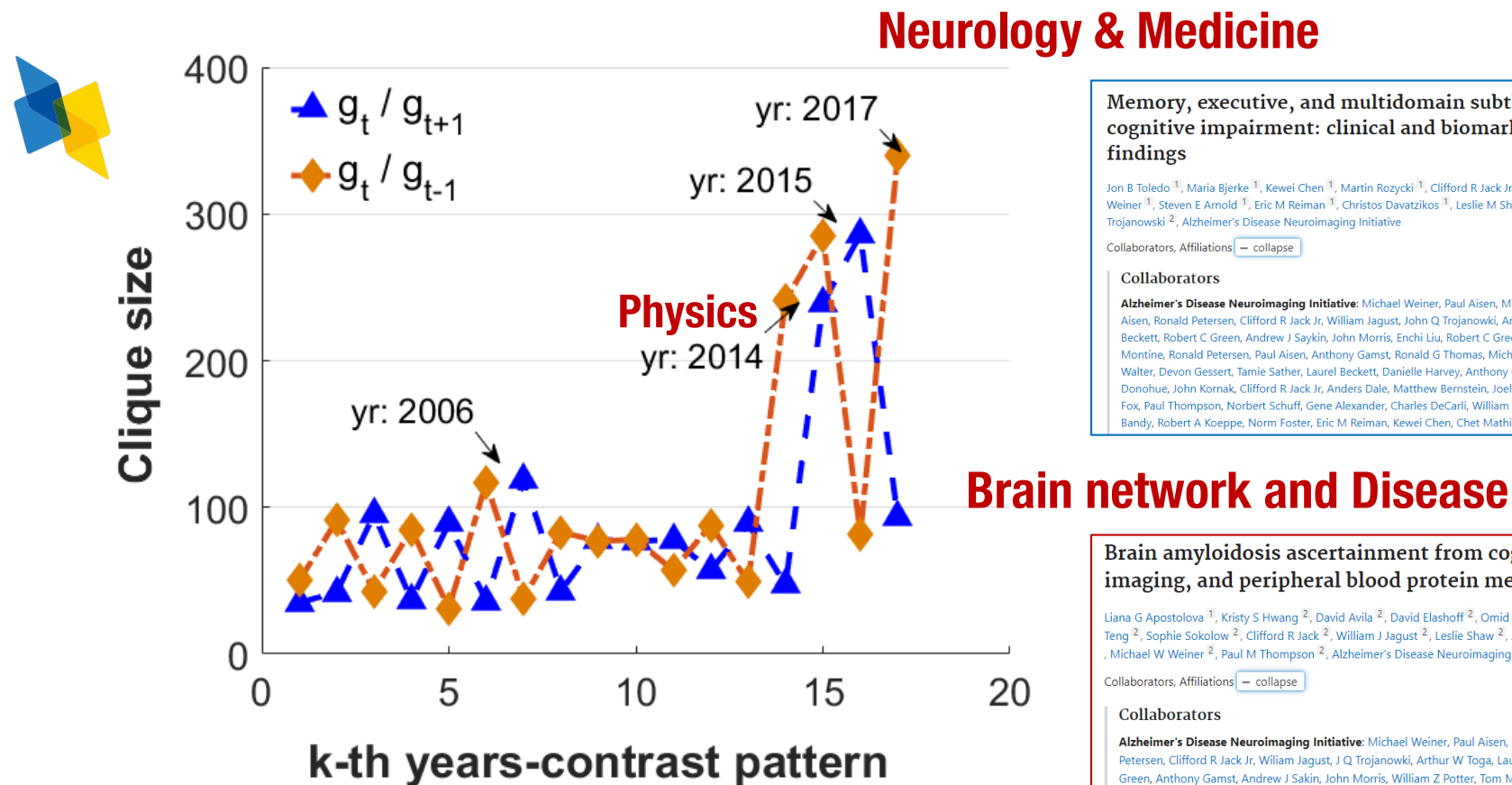
Exp2. Parameter Test

- Optimal graph spectral for the densest subgraph



Exp3. Effectiveness Test

- Contrast dense subgraphs in dynamic graph



Memory, executive, and multidomain subtle cognitive impairment: clinical and biomarker findings

Jon B Toledo¹, Maria Bjerke¹, Kewei Chen¹, Martin Rozyski¹, Clifford R Jack Jr¹, Michael W Weiner¹, Steven E Arnold¹, Eric M Reiman¹, Christos Davatzikos¹, Leslie M Shaw¹, John Q Trojanowski², Alzheimer's Disease Neuroimaging Initiative

Collaborators, Affiliations [collapse](#)

Collaborators

Alzheimer's Disease Neuroimaging Initiative: Michael Weiner, Paul Aisen, Michael Weiner, Paul Aisen, Ronald Petersen, Clifford R Jack Jr, William Jagust, John Q Trojanowski, Arthur W Toga, Laurel Beckett, Robert C Green, Andrew J Saykin, John Morris, Enchi Liu, Robert C Green, Tom Montine, Ronald Petersen, Paul Aisen, Anthony Gamst, Ronald G Thomas, Michael Donohue, Sarah Walter, Devon Gessert, Tamie Sather, Laurel Beckett, Danielle Harvey, Anthony Gamst, Michael Donohue, John Kornak, Clifford R Jack Jr, Anders Dale, Matthew Bernstein, Joel Felmlee, Nick Fox, Paul Thompson, Norbert Schuff, Gene Alexander, Charles DeCarli, William Jagust, Dan Bandy, Robert A Koeppe, Norm Foster, Eric M Reiman, Kewei Chen, Chet Mathis, John Morris, Nigel

Brain network and Disease

Brain amyloidosis ascertainment from cognitive, imaging, and peripheral blood protein measures

Liana G Apostolova¹, Kristy S Hwang², David Avila², David Elashoff², Omid Kohannim², Edmond Teng², Sophie Sokolow², Clifford R Jack², William J Jagust², Leslie Shaw², John Q Trojanowski², Michael W Weiner², Paul M Thompson², Alzheimer's Disease Neuroimaging Initiative

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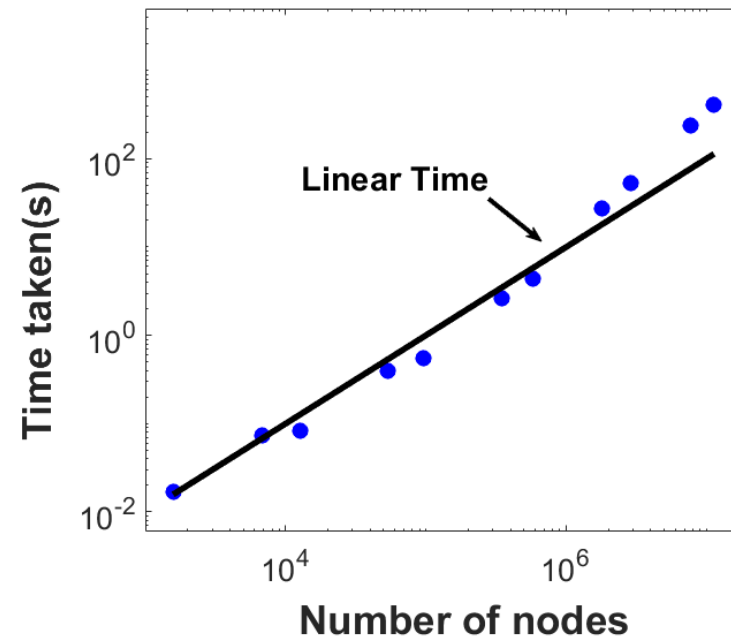
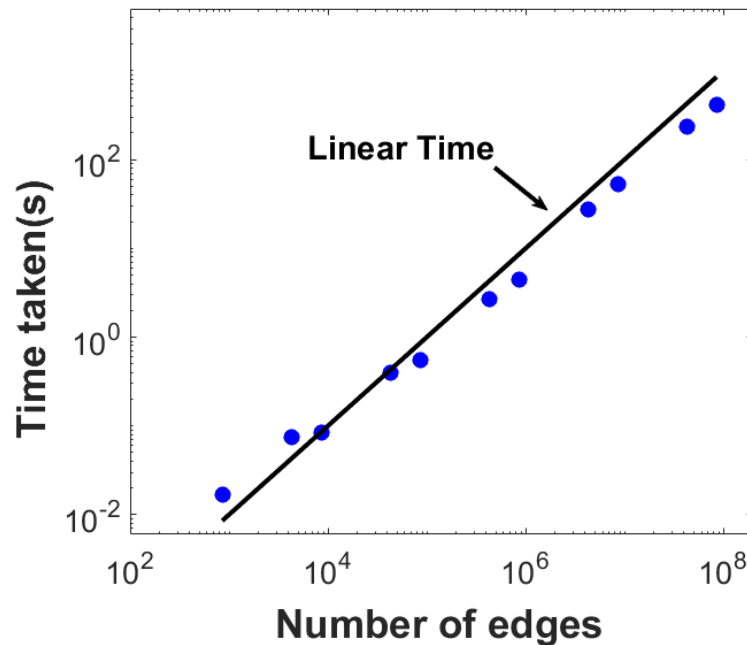
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2017 yr, 300+ co-authors

Exp4. Scalability Test

- Time complexity: $O(K \cdot |E| + K \cdot |E(\tilde{S})| \log|\tilde{S}|)^1$
- Linear scalable with the size of the graph



¹ \tilde{S} : max-size candidate S_t , $|\tilde{S}| \ll |S|$



Experiments

Properties of **SPECGREEDY**:

- ✓ **Faster**: detects the high-quality the densest subgraphs faster than the Greedy algo.
- ✓ **Effective**: detects the contrast dense subgraph pattern in scholar co-authorship
- ✓ **Scalable**: runs linearly w.r.t graph size

Road Map

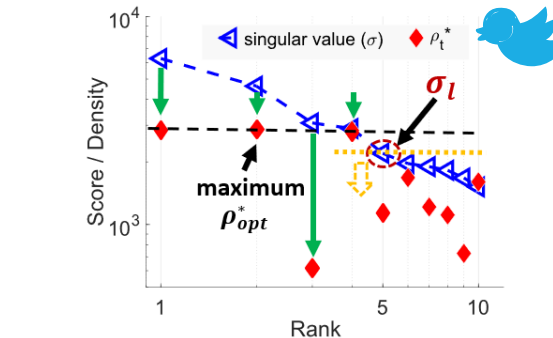
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Summary

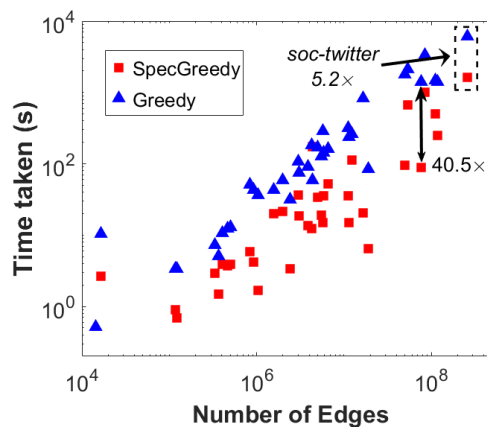
• SPECGREEDY: Unified Dense Subgraph Detection

Method	matrix P	matrix Q
MinQuotientCut [9]	$\mathbf{A} \quad - \quad \mathbf{D} = -\mathbf{L}$	\mathbf{I}
Charikar [6]	\mathbf{A}	\mathbf{I}
Fraudar [14]	$\mathbf{A} + 2 \mathbf{D}_w$	\mathbf{I}
SPARSECUTDS ¹ [24]	$\mathbf{A} - \frac{2 \cdot \alpha}{2\alpha+1} \mathbf{D}$	\mathbf{I}
TEMPDS [33]	\mathbf{A}_t	$\mathbf{A}_{t-1} + 2 \mathbf{I} = \tilde{\mathbf{A}}_{t-1}$
Risk-averse DS [30]	$\mathbf{A}^+ + \lambda_1 \mathbf{I} = \tilde{\mathbf{A}}^+$	$\mathbf{A}^- + \lambda_2 \mathbf{I} = \tilde{\mathbf{A}}^-$
GENDS²	$\mathbf{A} + 2 \mathbf{D}_c$	$\mathbf{A}' + \gamma \mathbf{I} = \tilde{\mathbf{A}}'$

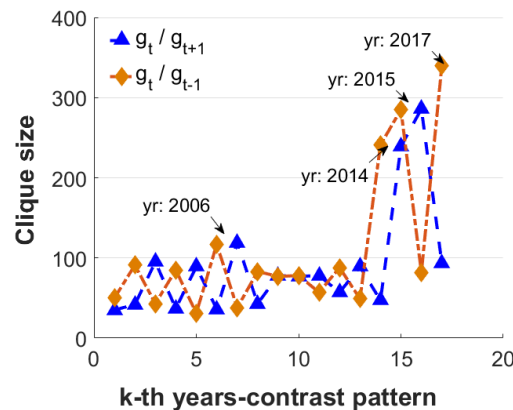
Unified formulation (GenDS)



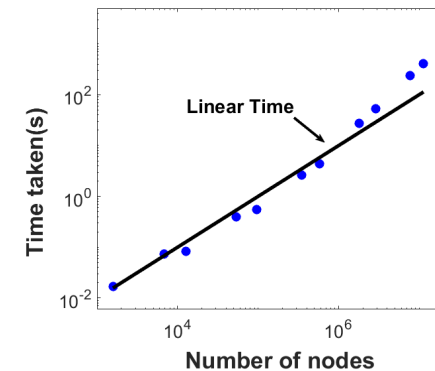
Graph spectral guaranty



☒ **Faster**



☒ **Effective**



☒ **Scalable**



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INSTITUTE OF COMPUTING TECHNOLOGY, CHINESE ACADEMY OF SCIENCES



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code



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